



Lecture 3

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ECE 336

***Stochastic Models for the Analysis of Computer Systems
and Communication Networks***



Outline of Lecture 3

- Summary of Lecture 2
- Properties of Exponential and Geometric Distributions
- Properties of the Poisson Process
- Steady-state solutions to DTMC and CTMC
- State Transition Rate Diagrams
- Applications to $M/M/1$, $M/M/1/N$, $M/M/m$, $M/M/\infty$, $M/M/m/m$ Queues
- Performance measures for $M/M/1$, $M/M/\infty$, and $M/M/m$ Queues

Summary of Lecture 2

DTMC:

- Specified by

$$P(X_{n+1} = j / X_n = i) = P_{ij} \begin{cases} \rightarrow \text{function of } n \\ \rightarrow \text{Independent of } n \end{cases}$$

non-homogeneous

homogeneous

- Unconditional probabilities

$$\underline{p}(n+1) = P^T \underline{p}(n); \sum_{j=0}^N P_{ij} = 1 \quad \forall i = 0, 1, \dots, N$$

In the steady state, $\underline{p} = P^T \underline{p}$ & $\lim_{n \rightarrow \infty} (P^T)^n = [\underline{p} \quad \underline{p} \quad \dots \quad \underline{p}]$; also $\sum_{i=0}^N p_i = 1$

The pmf of # of stages spent in a given state i : $P(\tau_i = n) = (1 - P_{ii}) P_{ii}^{n-1}$
Geometric pmf

CTMC:

- Specified by

$$P(X(t + \Delta t) = j / X(t) = i) = q_{ij}(t) \Delta t + o(\Delta t), i \neq j; P(X(t + \Delta t) = i / X(t) = i) = 1 + q_{ii}(t) \Delta t + o(\Delta t)$$

$$\sum_{j=0}^N q_{ij}(t) = 0 \Rightarrow q_{ii}(t) < 0 \text{ since } q_{ij}(t) \geq 0 \Rightarrow Q \underline{e} = \underline{0}$$

$$\dot{\underline{p}} = Q^T \underline{p}, p_i(t) = \text{Prob}\{X(t) = i\}; Q(t) \text{ time invariant} \Rightarrow \text{homogeneous Markov chain}$$

$$\text{In the steady state, } \dot{\underline{p}} = \underline{0} = Q^T \underline{p}$$

$$\text{The density of time spent in a given state } i: h_{\tau_i}(t) = \lambda_i \exp(-\lambda_i t), \lambda_i = -q_{ii}$$



Issues Addressed in Lecture 3

□ Questions:

- What are the properties of exponential density & Geometric pmf ?
- Since inter arrival times of a Poisson distribution are exponentially distributed, what are the properties of Poisson process ?
- Is there a general way of looking at arrival and departure processes?
.... Renewal processes
- When does steady-state solutions exist for DTMC and CTMC?
- Are there simple ways of visualizing state transitions?
- How to analyze simple queues ?



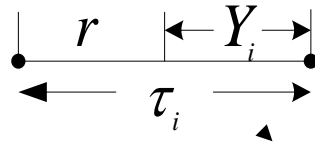
Exponential Density and Geometric pmf -1

□ Properties of Exponential pdf and Geometric pmf

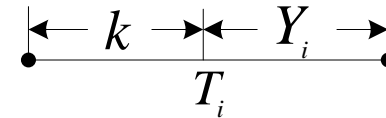
a) Memory-less Property:

Suppose that we have already spent in state i for r time units (k steps).
want:

CTMC



DTMC



$$\begin{aligned}
 P(Y_i > t | \tau_i > r) &= P(\tau_i > r + t | \tau_i > r) \\
 &= \frac{P(\tau_i > r + t, \tau_i > r)}{P(\tau_i > r)} \\
 &= \frac{P(\tau_i > r + t)}{P(\tau_i > r)} \\
 &= e^{-\lambda t} \\
 &= P(\tau_i > t)
 \end{aligned}$$

Unconditional pdf

$$\begin{aligned}
 P(Y_i > n | T_i > k) \quad n = 1, 2, \dots \\
 &= P(T_i > n + k | T_i > k) \\
 &= \frac{P(T_i > n + k, T_i > k)}{P(T_i > k)} \\
 &= \frac{P(T_i > n + k)}{P(T_i > k)} = P_{ii}^n \\
 &= P(T_i > n)
 \end{aligned}$$

Unconditional pmf



Exponential Density and Geometric pmf -2

Implication: Suppose τ denotes the time to failure of a component. Result says that if the component has lasted until time r , the distribution of the remaining (residual) time $Y = \tau - r$ is the same as the time to failure of a new component. If the time to failure is exponentially distributed, then a used component is “as good as new”.

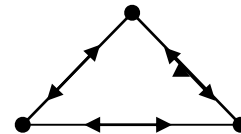
b) Hazard Rate:

$$Z_i(t) = \frac{h_{\tau_i}(t)}{1 - H_{\tau_i}(t)} = \frac{\lambda_i e^{-\lambda_i t}}{e^{-\lambda_i t}} = \lambda_i; \quad Z_n = \frac{P(T_i = n)}{1 - P(T_i \geq n)} = \frac{P(T_i = n)}{1 - P(T_i > n - 1)}$$

$$= 1 - P_{ii}$$

Hazard rate is a constant

Markovian property (e.g., Poisson)



Exponential density

Constant hazard rate



Exponential Density and Geometric pmf - 3

c) Among all the densities having the same mean, exponential density has the largest entropy.

$$H = -E\{\ln h(\tau)\} = -\int_0^{\infty} h(\tau) \ln h(\tau) d\tau$$

Define Lagrangian :

$$L = -\int_0^{\infty} h(\tau) \ln h(\tau) d\tau - \gamma \left[\int_0^{\infty} h(\tau) d\tau - 1 \right] - \beta \left[\int_0^{\infty} \tau h(\tau) d\tau - \bar{\tau} \right]$$

$$-\ln h(\tau) - (1 + \gamma) - \beta\tau = 0 \Rightarrow h(\tau) = e^{-(1+\gamma+\beta\tau)}$$

$$\text{From normalization \& mean constraints: } \frac{e^{-(1+\gamma)}}{\beta} = 1 \Rightarrow e^{-(1+\gamma)} = \beta \Rightarrow h(\tau) = \beta e^{-\beta\tau}$$

$$\text{Mean constraint yields: } \int_0^{\infty} \beta\tau e^{-\beta\tau} d\tau = \bar{\tau} = \frac{1}{\beta} \Rightarrow \beta = \frac{1}{\bar{\tau}}. \text{ so, } h(\tau) = \frac{1}{\bar{\tau}} e^{-\tau/\bar{\tau}}$$

$$\text{Can show similar result for geometric pmf : } H = -\sum_{n=0}^{\infty} P(\tau_i = n) \ln P(\tau_i = n)$$



Exponential Density and Geometric pmf - 4

d) Suppose have n independent random variables $\tau_1, \tau_2, \dots, \tau_m$ with parameters $\lambda_1, \lambda_2, \dots, \lambda_m$, respectively. Then

$$\tau = \min(\tau_1, \tau_2, \dots, \tau_m) \text{ is exponential with parameter } \sum_{i=1}^m \lambda_i$$

Proof:

$$P(\tau \geq t) = P(\tau_1 \geq t)P(\tau_2 \geq t) \dots P(\tau_m \geq t) = e^{-\left(\sum_{i=1}^m \lambda_i\right)t}$$

$$\Rightarrow \text{exponential with parameter } \lambda = \sum_{i=1}^m \lambda_i$$

Similarly for geometric pmf :

$$T = \min(T_1, T_2, \dots, T_m) \text{ where } T_i = \text{Geometric with parameter } P_{ii}$$

$$\text{Then, } P(T > n) = P(T_1 > n)P(T_2 > n) \dots P(T_m > n) = \left(\prod_{i=1}^m P_{ii}^n\right) = \left(\prod_{i=1}^m P_{ii}\right)^n$$

$$\therefore T \text{ is geometric with parameter } P = \left(\prod_{i=1}^m P_{ii}\right) \Rightarrow P(T = n) = (1 - P)P^{n-1}$$

$$\text{Aside: } P(\tau_i < \tau_j) = \int_0^{\infty} P(\tau_i < \tau_j | \tau_i = t) h_{\tau_i}(t) dt = \lambda_i \int_0^{\infty} e^{-\lambda_j t} e^{-\lambda_i t} dt$$

$$= \frac{\lambda_i}{\lambda_i + \lambda_j} \Rightarrow \text{relative rate}$$



Exponential Density and Geometric pmf - 5

e) Suppose have n independent random variables $\tau_1, \tau_2, \dots, \tau_m$ with the same parameter λ . Then

$$S_m = \sum_{i=1}^m \tau_i \text{ is Erlang-}m \text{ (or Gamma}(m, \lambda)) \text{ distribution}$$

Proof:

$$\text{Recall MGF : } L_{\tau_i}(s) = \frac{\lambda}{(s + \lambda)}$$

$$\Rightarrow L_{S_m}(s) = \left(\frac{\lambda}{s + \lambda} \right)^m \Rightarrow h_{S_m}(t) = \frac{\lambda(\lambda t)^{m-1}}{(m-1)!} e^{-\lambda t}; t \geq 0$$

$$\Rightarrow E(S_m) = \frac{m}{\lambda}; E(S_m^2) = \frac{m(m+1)}{\lambda^2}$$

$$\Rightarrow \sigma_{S_m}^2 = \frac{m}{\lambda^2} \Rightarrow C_{S_m} = \frac{1}{\sqrt{m}} \rightarrow 0 \text{ as } m \rightarrow \infty$$

Application:

Pdf of the time until m^{th} arrival.

$X(t)$ Poisson $\Leftrightarrow X(t) = \max\{m: S_m \leq t\}$ & $S_0 = 0$.

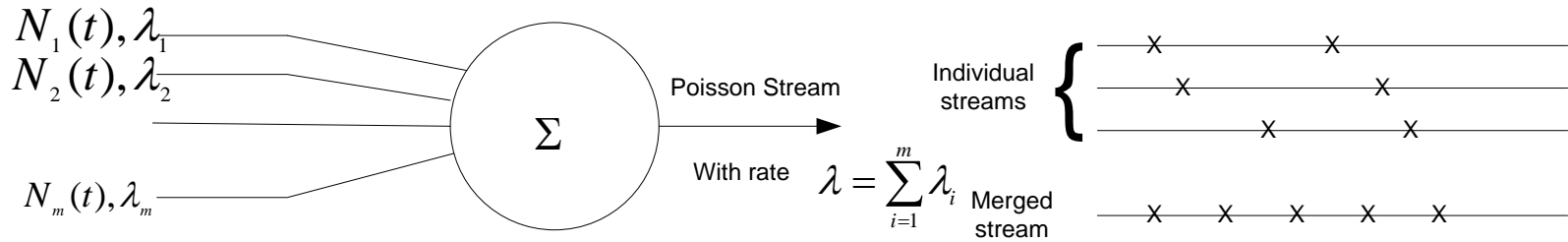
Aside: Suppose $\tau_1, \tau_2, \dots, \tau_m$ have different parameters, $\lambda_1, \lambda_2, \dots, \lambda_m; \lambda_i \neq \lambda_j$

$$\Rightarrow L_{S_m}(s) = \prod_{i=1}^m \frac{\lambda_i}{s + \lambda_i} \Rightarrow h_{S_m}(t) = \sum_{i=1}^m \lambda_i e^{-\lambda_i t} \prod_{\substack{j=1 \\ j \neq i}}^m \frac{\lambda_j}{\lambda_j - \lambda_i}; t \geq 0$$



Properties of Poisson Process - 1

a) Merging of Poisson Processes results in a Poisson Process



Proof:

$$N(t) = N_1(t) + \dots + N_m(t)$$

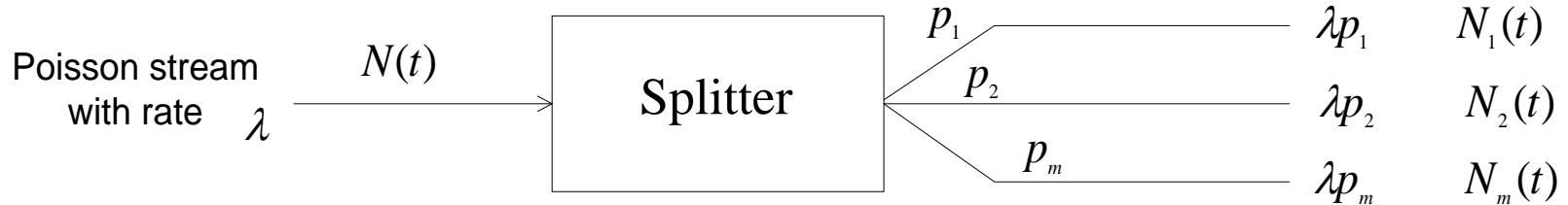
sum of i.i.d. \Rightarrow convolution \Rightarrow product of MGFs

$$G_N(z, t) = e^{(\sum_{i=1}^m \lambda_i)t(z-1)} = e^{\lambda t(z-1)} \Rightarrow p_n(t) = \frac{(\lambda t)^n}{n!} e^{-\lambda t}; n = 0, 1, 2, \dots$$



Properties of Poisson Process - 2

b) Splitting or Decomposition of a Poisson Process



$$N(t) = N_1(t) + N_2(t) + \dots + N_m(t)$$

$$P(N_1, N_2, \dots, N_m | N) = \frac{N!}{N_1! N_2! \dots N_m!} p_1^{N_1} p_2^{N_2} \dots p_m^{N_m}$$

$$P(N_1, N_2, \dots, N_m; N) = P(N_1, N_2, \dots, N_m | N) P(N)$$

$$= e^{(-\sum \lambda p_i)t} N! \frac{p_1^{N_1} p_2^{N_2} \dots p_m^{N_m} (\lambda t)^{N_1 + N_2 + \dots + N_m}}{N_1! N_2! \dots N_m! N!}$$

$$= \prod_{i=1}^m \frac{e^{-\lambda p_i t} (\lambda p_i t)^{N_i}}{N_i!}$$

\Rightarrow Individual streams are independent and poisson with rates, λp_i



Properties of Poisson Process - 3

c) Conditional Distribution of the Arrival Times

Suppose one event of a Poisson process has taken place by time, t . Given this information, the distribution of the time at which the event occurred is *uniform* in $[0, t]$.

$$\begin{aligned} P(\tau_1 < s \mid N(t) = 1) &= \frac{P(\tau_1 < s, N(t) = 1)}{P(N(t) = 1)} \\ &= \frac{P\{1 \text{ event in } [0, s], 0 \text{ events in } [s, t]\}}{P(N(t) = 1)} \\ &= \frac{P\{1 \text{ event in } [0, s]\}P\{0 \text{ events in } [s, t]\}}{P(N(t) = 1)} \\ &= \frac{\lambda s e^{-\lambda s} e^{-\lambda(t-s)}}{\lambda t e^{-\lambda t}} = \frac{s}{t} \Rightarrow \text{uniform} \end{aligned}$$

Aside: Given that $N(t) = n$, the n arrival times S_1, S_2, \dots, S_n are i.i.d. and uniformly distributed in the interval $[0, t]$

Inter - arrival times : $\tau_i = S_i - S_{i-1}; S_0 = 0$

$$\begin{aligned} P(s_1, s_2, \dots, s_n \mid n) &= \frac{h(s_1, s_2, \dots, s_n, n)}{P(N(t) = n)} = \frac{\lambda e^{-\lambda s_1} \lambda e^{-\lambda(s_2 - s_1)} \dots e^{-\lambda(s_n - s_{n-1})} e^{-\lambda(t - s_n)}}{e^{-\lambda t} (\lambda t)^n / n!} \\ &= \frac{n!}{t^n}; 0 < s_1 < \dots < s_n < t \end{aligned}$$



Properties of Poisson Process - 4

d) Poisson pmf is a Special Case of Binomial pmf

$$p(n) = \binom{N}{n} p^n (1-p)^{N-n}$$

$$\text{Let } p = \lambda \Delta t$$

$$N = \frac{t}{\Delta t}$$

$$\lim_{\Delta t \rightarrow 0} Np = \lambda t$$

$$\Rightarrow p(n) = \frac{N!}{N-n!n!} (\lambda \Delta t)^n (1 - \lambda \Delta t)^{N-n} = \frac{N!}{N-n!n!} \frac{(\lambda t)^n}{N^n} \left(1 - \frac{\lambda t}{N}\right)^{N-n}$$

$$= \frac{N(N-1)\dots(N-n+1)}{n!N^n} (\lambda t)^n \left(1 - \frac{\lambda t}{N}\right)^{\frac{N}{\lambda t} \cdot \frac{\lambda t(N-n)}{N}}$$

$$= \frac{N(N-1)\dots(N-n+1)}{n!N^n} (\lambda t)^n e^{-\lambda t \frac{N-n}{N}}$$

$$\text{As } N \rightarrow \infty \quad p(n) = \frac{(\lambda t)^n}{n!} e^{-\lambda t}; \quad n = 0, 1, 2, \dots$$



Digression: Renewal Processes

- Suppose $\{\tau_i : i = 0, 1, 2, \dots\}$ are i.i.d. nonnegative random variables (*not necessarily exponential*). Let $S_n = \sum_{i=1}^n \tau_i$

$$S_{n+1} = S_n + \tau_{n+1} \Rightarrow h_{S_{n+1}}(t) = h_{S_n}(t) * h_{\tau}(t) = \int_0^t h_{S_n}(t - \sigma) h_{\tau}(\sigma) d\sigma$$

$X(t) \sim$ counting process

$$P(X(t) = n) = P\{S_n \leq t < S_{n+1}\} = P(S_n \leq t) - P(S_{n+1} \leq t) = H_{S_n}(t) - H_{S_{n+1}}(t)$$

$$\Rightarrow P(X(t) \geq n) = H_{S_n}(t) \Rightarrow M(t) = E[X(t)] = \sum_{n=1}^{\infty} H_{S_n}(t) = H_{\tau}(t) + \sum_{n=1}^{\infty} H_{S_{n+1}}(t)$$

$$\Rightarrow M(t) = H_{\tau}(t) + \sum_{n=1}^{\infty} \int_0^t H_{S_n}(t - \sigma) h_{\tau}(\sigma) d\sigma = H_{\tau}(t) + \int_0^t \sum_{n=1}^{\infty} H_{S_n}(t - \sigma) h_{\tau}(\sigma) d\sigma$$

$$\Rightarrow M(t) = H_{\tau}(t) + \int_0^t M(t - \sigma) h_{\tau}(\sigma) d\sigma \quad \text{Fundamental Renewal Equation}$$

$$\text{Let Rate of renewals, } m(t) = \frac{dM(t)}{dt} = \sum_{n=1}^{\infty} h_{S_n}(t) \Rightarrow m(t) = h_{\tau}(t) + \int_0^t m(t - \sigma) h_{\tau}(\sigma) d\sigma \quad \text{Renewal Equation}$$

$$\text{In Laplace domain : } L_m(s) = \frac{L_{\tau}(s)}{1 - L_{\tau}(s)} \Rightarrow \lim_{t \rightarrow \infty} m(t) = \lim_{s \rightarrow 0} s L_m(s)$$

$$\lim_{t \rightarrow \infty} m(t) = \frac{L_{\tau}(s) + s \frac{dL_{\tau}(s)}{ds}}{-\frac{dL_{\tau}(s)}{ds}} \Big|_{s=0} = \frac{1}{E[\tau]} \Rightarrow \text{long term rate of renewals} = \frac{1}{\text{inter-renewal time}}$$



Poisson Process is a Renewal Process

- Suppose $\{\tau_i : i = 0, 1, 2, \dots\}$ are i.i.d. exponential random variables

$$\text{Let } S_n = \sum_{i=1}^n \tau_i$$

$$\text{Exponential} \Rightarrow L_\tau(s) = \frac{\lambda}{(s + \lambda)}$$

$$\Rightarrow L_m(s) = \frac{\lambda}{s} \Rightarrow m(t) = \lambda$$

$\Rightarrow M(t) = \lambda t = \text{Expected number of renewals in } [0, t] = \text{mean of Poisson}$

$$\text{Also, } L_{S_n}(s) = [L_\tau(s)]^n \Rightarrow h_{S_n}(t) = \text{Erlang}(n, \lambda) = \frac{\lambda(\lambda t)^{n-1} e^{-\lambda t}}{(n-1)!}$$

$$\Rightarrow H_{S_n}(t) = \int_0^t h_{S_n}(\sigma) d\sigma = 1 - \sum_{k=0}^{n-1} \frac{(\lambda t)^k}{k!} e^{-\lambda t}$$

$$\Rightarrow P(X(t) = n) = H_{S_n}(t) - H_{S_{n+1}}(t) = \frac{(\lambda t)^n e^{-\lambda t}}{n!}; n = 0, 1, 2, \dots$$



Steady-state Probabilities -1

- When does steady-state solution to $\underline{p} = P^T \underline{p}$ & $\underline{0} = Q^T \underline{p}$ exist?

We need n -step transition probabilities. Let

$$\Phi_{ij}(n) = P\{X_{m+n} = j / X_m = i\} = P\{X_n = j / X_0 = i\} \text{ time-homogeneous MC}$$

$$\begin{aligned}\Phi_{ij}(n) &= \sum_{k=0}^N P\{X_n = j / X_{n-1} = k\} \cdot P\{X_{n-1} = k / X_0 = i\} \\ &= \sum_{k=0}^N \Phi_{ik}(n-1) P_{kj}\end{aligned}$$

$$\Phi(n) = \Phi(n-1)P \Rightarrow \Phi(n) = P^n \text{ Similar to state-transition matrix in control theory}$$

- Two states i and j communicate if $\exists n, n' \ni \Phi_{ij}(n) = (P^n)_{ij} > 0$ and $\Phi_{ji}(n') = (P^{n'})_{ji} > 0$
- state i communicates with itself
 - if i and j communicate, then j and i communicate
 - if i and j communicate, and j and k communicate, then i and k communicate (transitive)
 - communicating states define a *class* \Rightarrow comm. divides states into *classes*
- If all states communicate (i.e., one class) \Rightarrow Markov chain is *irreducible*



First Passage Probabilities

First Passage Probabilities and Recurrence Times

$$f_{ij}(n) = P\{\text{first return to state } j \text{ occurs after } n \text{ steps of leaving state } i\}$$

$$= P\{\text{first time } X(n) = j \mid X(0) = i\}$$

\Rightarrow State j can be reached for the first time at the k^{th} step with probability $f_{ij}(k)$ and again in the remaining $n - k$ steps with

$$\text{probability } \Phi_{jj}(n - k) \Rightarrow \Phi_{ij}(n) = \sum_{k=1}^n f_{ij}(k) \Phi_{jj}(n - k); n \geq 1$$

Initial conditions: $f_{ij}(0) = 0$; $f_{ij}(1) = P_{ij}$; $\Phi_{jj}(0) = 1$, $\Phi_{ij}(0) = 0$.

$$\Rightarrow \text{For } i = j, \Phi_{ii}(z) = 1 + F_{ii}(z)\Phi_{ii}(z) \Rightarrow \Phi_{ii}(z) = \frac{1}{1 - F_{ii}(z)}$$

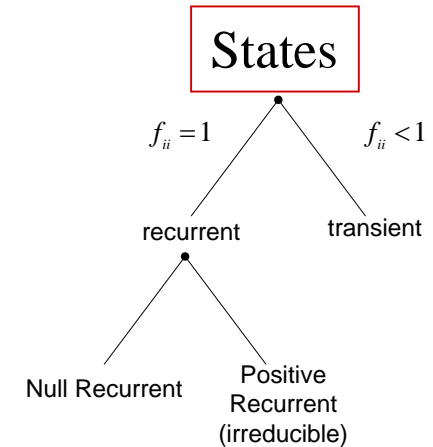
$$f_{ii} = \sum_{n=1}^{\infty} f_{ii}(n) = \text{Prob}\{\text{ever returning to state } i\}$$

$f_{ii} = 1 \Rightarrow \text{persistent (recurrent)}$; $f_{ii} < 1 \Rightarrow \text{transient}$

Mean Recurrence Time:

$$M_i = \sum_{n=1}^{\infty} n f_{ii}(n) \quad M_i = \infty \text{ recurrent null}$$

$$M_i < \infty \text{ recurrent nonnull (Positive recurrent)}$$

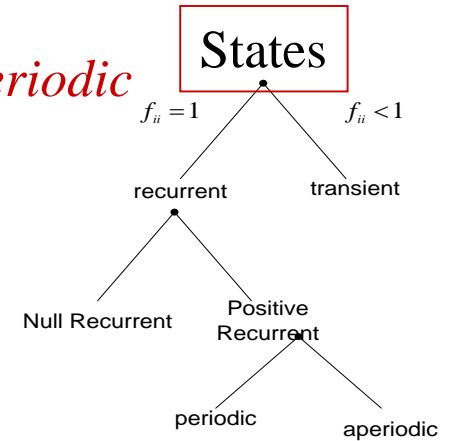




Steady-state Probabilities - 2

Further classification of Markov-chain states

- If the return to state j occurs only at $\gamma, 2\gamma, 3\gamma, \dots$ ($\gamma > 1$), $\{\Phi_{ii}(n) = 0 \text{ unless } n = k\gamma\}$, then state j is *periodic*
- If $\gamma = 1$ for state i , $\{\Phi_{ii}(n) > 0 \forall n\}$ then it is said to be *aperiodic*



If one state in a class is *aperiodic*, all states in that class are
 If one state in a class is *transient*, all states in that class are
 If one state in a class is *recurrent*, all states in that class are
 If one state in a class is *periodic*, all states in that class are

Irreducible \Rightarrow recurrent non-null (positive recurrent) and aperiodic

Irreducible aperiodic (i.e., ergodic) Markov chains have:

1) $p_i = 0 \quad \forall i \geq 0 \Rightarrow$ no steady - state distribution (e.g., arrival rate $>$ service capacity)

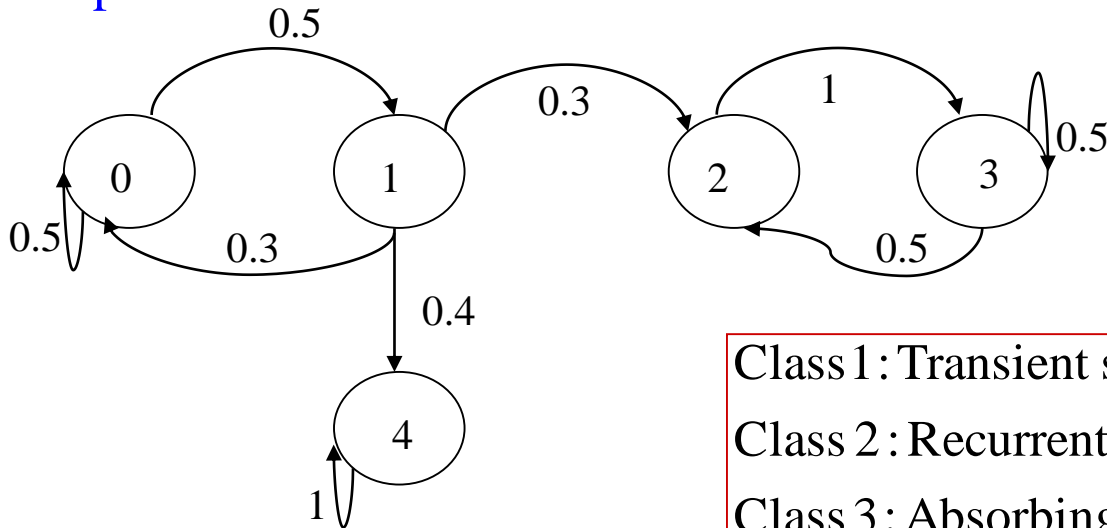
2) $p_i > 0 \quad \forall i \Rightarrow \{p_i / i \geq 0\} = p_i = \frac{1}{M_i}$ is the unique stationary distribution

Absorbing state: $P_{ii} = 1 \Rightarrow M_i = 1 \Rightarrow p_i = 1$



Examples of State Classification - 1

Examples:



- Class 1: Transient states (0,1)
- Class 2: Recurrent states (2,3)
- Class 3: Absorbing state (4)

Steady state probabilities can be one of these two :

Two unity eigen values since two recurrent classes

- $[0\ 0\ 0\ 0\ 1] \Rightarrow$ chain is trapped in state 4
- $[0\ 0\ 0.25\ 2\ 2\ 0.50\ 4\ 4\ 0.24\ 3\ 4]$
 \Rightarrow finite probability of getting trapped in class of states (2,3)



Examples of State Classification - 2

Examples:

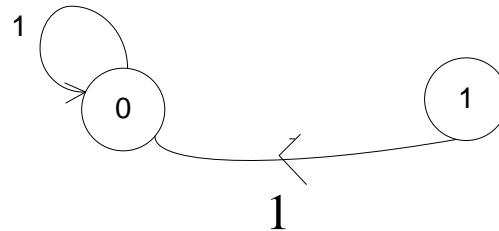
$$P = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \Rightarrow \Phi(n) = P^n = \begin{cases} \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix} & n \text{ even} \\ \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} & n \text{ odd} \end{cases}$$

starting in state 0, we return to state 0
in even # of steps
 $\Rightarrow \gamma = 2$ periodic

$$P = \begin{pmatrix} 1/3 & 2/3 \\ 1/2 & 1/2 \end{pmatrix} \Rightarrow \{\Phi_{ii}(n) = (P^n)_{ii} > 0 \forall n\}$$

aperiodic Markov chain

$$\begin{bmatrix} 1 & 0 \\ 1 & 0 \end{bmatrix}$$



State 0: absorbing
State 1: transient

steady state probability: [1 0]

We will consider *irreducible, aperiodic Markov chains*



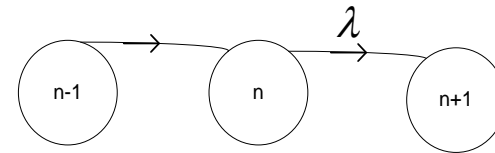
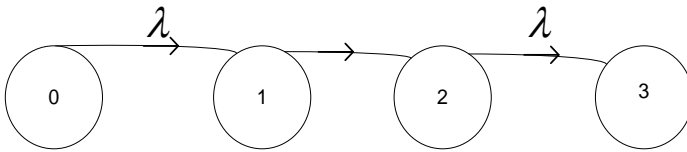
State Transition Rate Diagrams - 1

Examples

Poisson Process

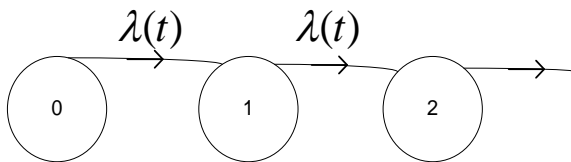
$$\dot{p}_n = \lambda p_{n-1}(t) - \lambda p_n(t)$$

$$\Rightarrow p_n(t) = \frac{(\lambda t)^n}{n!} e^{-\lambda t}$$



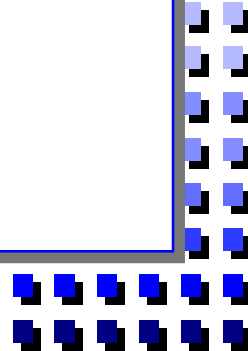
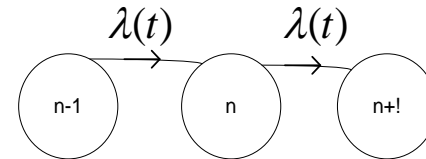
$$\begin{aligned} \dot{p}_i &= \sum_{j=0}^N q_{ji}(t) p_j(t) = \sum_{\substack{j \neq i \\ j=0}}^N q_{ji} p_j(t) + q_{ii}(t) p_i(t) \\ &= \underbrace{\sum_{\substack{j \neq i \\ j=0}}^N q_{ji} p_j(t)}_{\text{FLOW INTO NODE } i} - \underbrace{\sum_{\substack{j \neq i \\ j=0}}^N q_{ij} p_i(t)}_{\text{FLOW OUT OF NODE } i} \Rightarrow \text{No need to know } q_{ii} \end{aligned}$$

Non-homogeneous Poisson process



$$\dot{p}_n = \lambda(t) p_{n-1}(t) - \lambda(t) p_n(t)$$

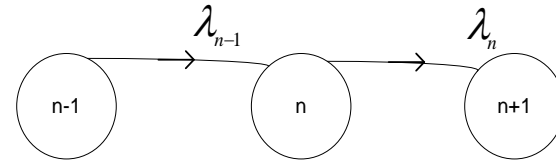
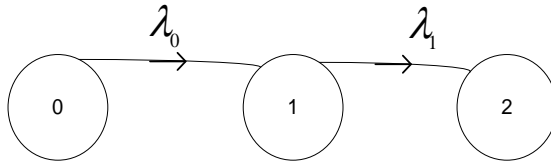
$$\Rightarrow p_n(t) = \frac{\left[\int_0^t \lambda(\tau) d\tau \right]^n}{n!} e^{-\int_0^t \lambda(\tau) d\tau}$$





State Transition Rate Diagrams - 2

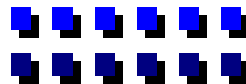
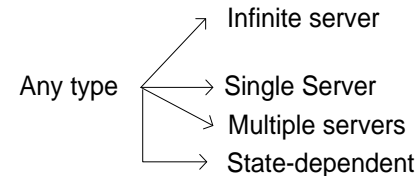
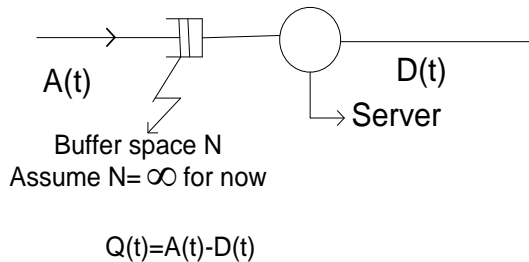
- State-dependent birth process



$$\dot{p}_n(t) = -\lambda_n p_n(t) + \lambda_{n-1} p_{n-1}(t)$$

$$\Rightarrow p_n(s) = \frac{\lambda_{n-1} p_{n-1}(s)}{s + \lambda_n} \Rightarrow p_n(t) = e^{-\lambda_n t} \left[\lambda_{n-1} \int_0^t p_{n-1}(\tau) e^{\lambda_n \tau} d\tau + p_n(0) \right]$$

- Birth-death processes: Forms the basis of *all* Markovian Queuing systems

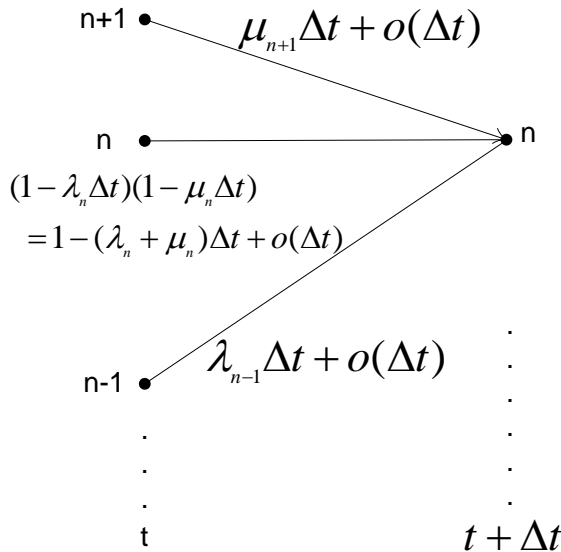




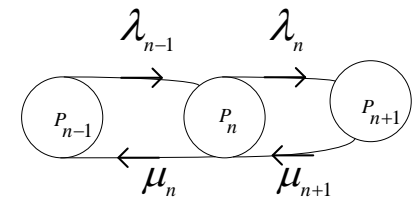
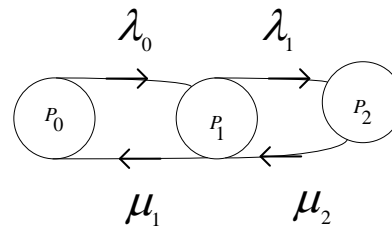
State Transition Rate Diagrams - 3

$$\left. \begin{array}{l} A(t) = \text{total births upto time } t \\ D(t) = \text{total deaths upto time } t \end{array} \right\} \Rightarrow Q(t) = \text{system length at time } t$$

We assume that in an interval of time Δt only one arrival or departure can occur



We assume $S=1$ (service demand=1)



$$\begin{aligned} \dot{p}_n(t) &= -(\lambda_n + \mu_n)p_n(t) + \lambda_{n-1}p_{n-1}(t) + \mu_{n+1}p_{n+1}(t); n \geq 1 \\ \dot{p}_0(t) &= -\lambda_0p_0(t) + \mu_1p_1(t) \end{aligned}$$



Application to Simple Queues - 1

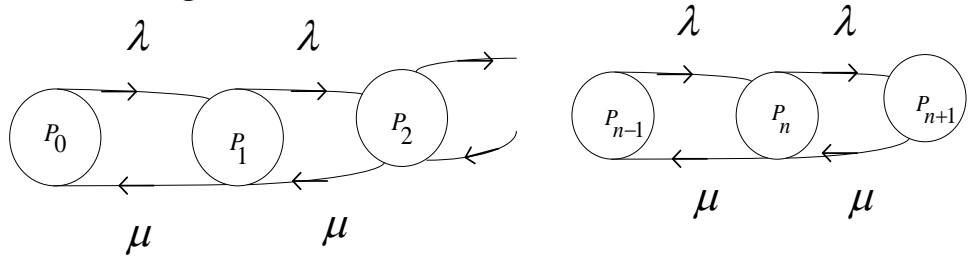
□ B-D processes model a wide variety of Markovian (exponential) queues

1) M/M/1 Queue $M \sim$ Poisson arrivals or exponential inter-arrival pdf
 $M \sim$ exponential service demands

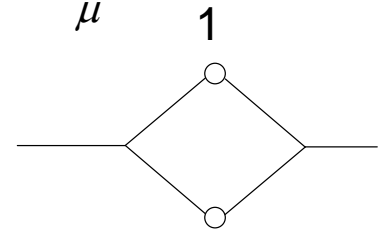


$1 \sim$ single server

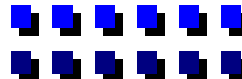
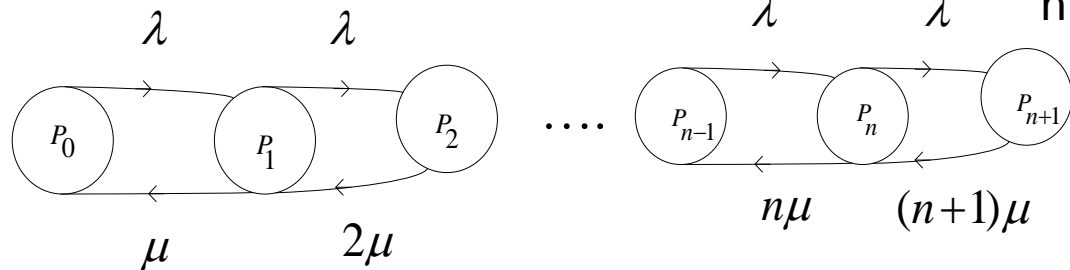
$$\left. \begin{matrix} \lambda_n = \lambda \\ \mu_n = \mu \end{matrix} \right\} \Rightarrow$$



2) M/M/∞ queue Infinite server case \Rightarrow No waiting



$$\left. \begin{matrix} \lambda_n = \lambda \\ \mu_n = n\mu \end{matrix} \right\} \Rightarrow$$



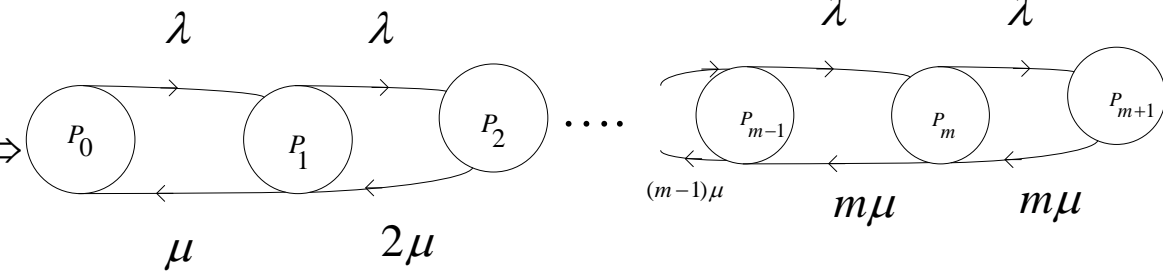


Application to Simple Queues - 2

3) M/M/m queue \Rightarrow m server case

$$\lambda_n = \lambda$$

$$\mu_n = \begin{cases} n\mu & n \leq m \\ m\mu & n > m \end{cases}$$

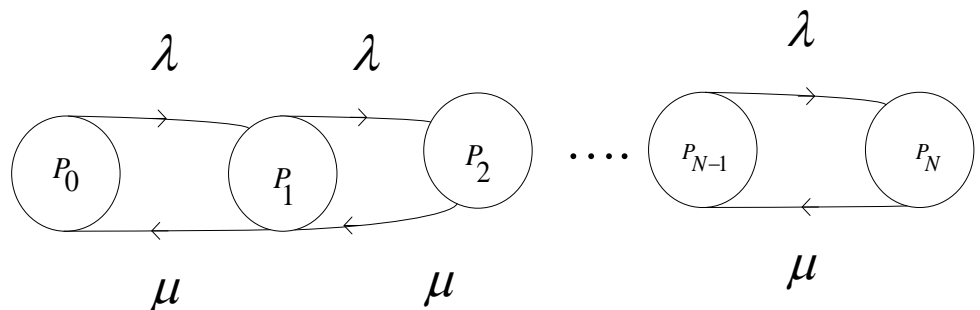


4) M/M/1/N queue \Rightarrow Model of a single machine with a finite buffer

At most N

$$\lambda_n = \begin{cases} \lambda & n \leq N - 1 \\ 0 & n \geq N \end{cases}$$

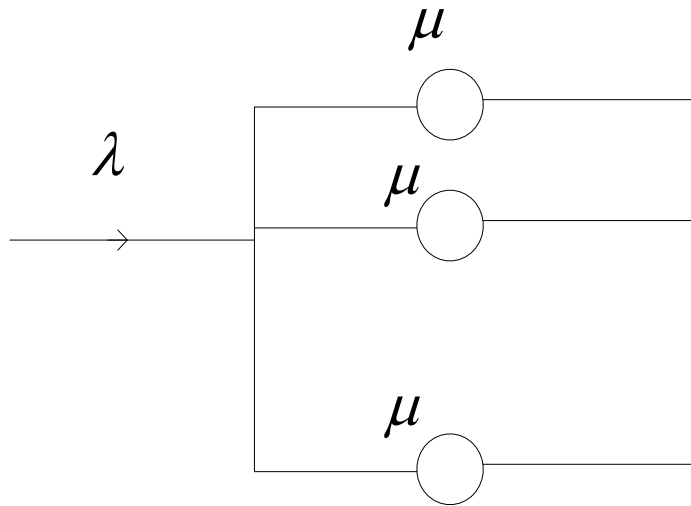
$$\mu_n = \begin{cases} \mu & n \leq N \\ 0 & n > N \end{cases}$$





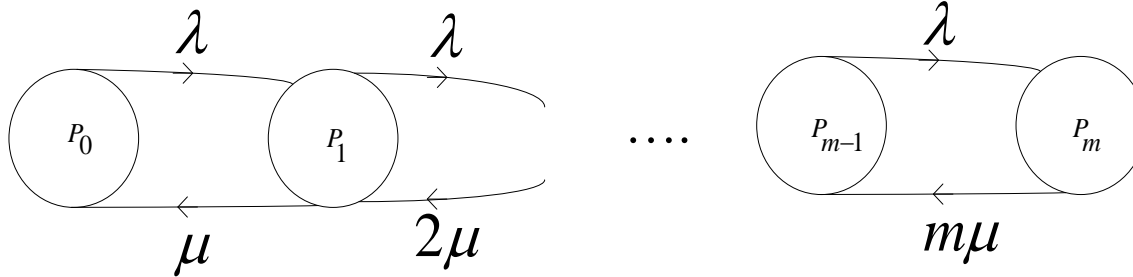
Application to Simple Queues - 3

5) M/M/m/m model \Rightarrow Blocked call loss system or Blocked calls cleared (BCC)



$$\lambda_n = \lambda \quad 0 \leq n \leq m-1$$

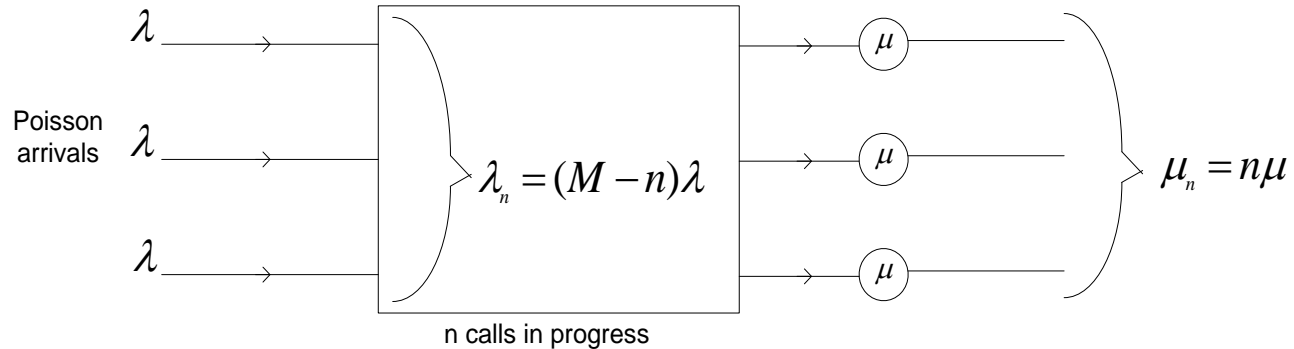
$$\mu_n = \begin{cases} n\mu & n \leq m \\ 0 & n > m \end{cases}$$





Application to Simple Queues - 4

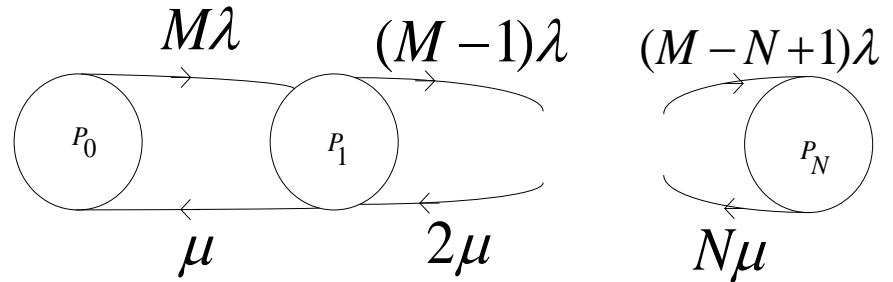
6) Engset model ($M > N$) $M/M/N/N/M$ queue



$$\lambda_n = (M - n)\lambda; \quad 0 \leq n \leq N$$

$$N \leq M$$

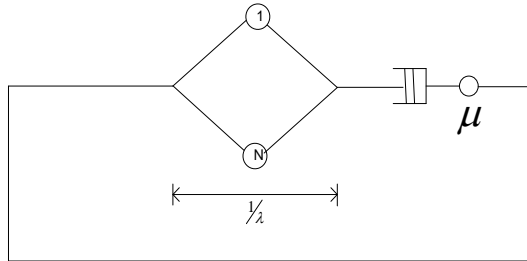
$$\mu_n = n\mu; \quad 1 \leq n \leq N$$





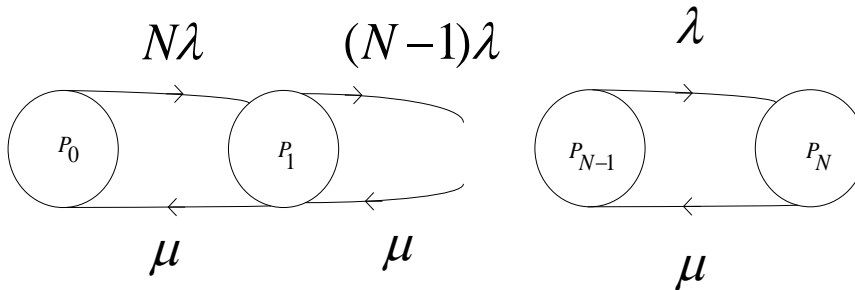
Application to Simple Queues - 5

7) Machine Repairman model...forms the basis of closed networks
M/M/1/N/N queue

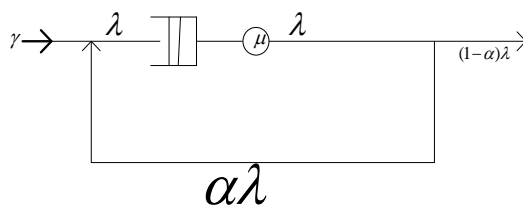


$$\lambda_n = \begin{cases} (N - n)\lambda & 0 \leq n \leq N \\ 0 & \text{otherwise} \end{cases}$$

$$\mu_n = \mu$$

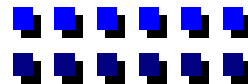


8) M/M/1 queue with feedback A simple open network



$$\Rightarrow \lambda_n = \frac{\gamma}{1 - \alpha}$$

$$\mu_n = \mu$$





Solution of Birth-Death Model - 1

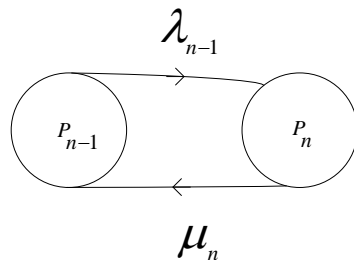
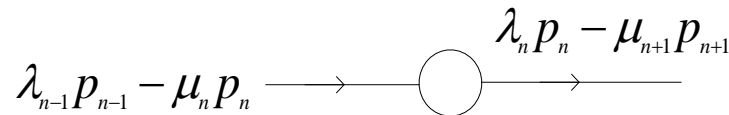
Coming back to birth-death model

$$p_n(t + \Delta t) = \lambda_{n-1}p_{n-1}(t)\Delta t + [1 - (\lambda_n + \mu_n)\Delta t]p_n(t) + \mu_{n+1}p_{n+1}(t)\Delta t$$

$$\left. \begin{aligned} \frac{dp_n(t)}{dt} &= \lambda_{n-1}p_{n-1}(t) - (\lambda_n + \mu_n)p_n(t) + \mu_{n+1}p_{n+1}(t); n \geq 1 \\ \frac{dp_0(t)}{dt} &= -\lambda_0p_0(t) + \mu_1p_1(t) \end{aligned} \right\} \text{known as global balance equations}$$

Solution via numerical methods. If $\lambda_n = \lambda, \mu_n = \mu$ (i.e., M/M/1 queue), solution involves Bessel functions. We will be content with steady - state solution only.

$$p_n = \lim_{t \rightarrow \infty} p_n(t)$$



$$\Rightarrow \lambda_{n-1}p_{n-1} = \mu_n p_n$$

Local balance
n detailed Balance Equations



Solution of Birth-Death Model - 2

$$\Rightarrow p_n = \frac{\lambda_{n-1}}{\mu_n} p_{n-1} = \frac{\lambda_0 \lambda_1 \dots \lambda_{n-1}}{\mu_1 \mu_2 \dots \mu_n} p_0 = \left[\prod_{i=0}^{n-1} \left(\frac{\lambda_i}{\mu_{i+1}} \right) \right] p_0$$

$$\text{Since } \sum_{n=0}^{\infty} p_n = 1 \Rightarrow p_0 = \left[1 + \sum_{n=1}^{\infty} \prod_{i=0}^{n-1} \frac{\lambda_i}{\mu_{i+1}} \right]^{-1}$$

- Once the distribution is known, we can obtain a variety of performance measures

1) Throughput: $X = \sum_{n=1}^{\infty} \mu_n p_n$

2) System (or Queue) length $Q = \sum_{n=1}^{\infty} n p_n$

3) Utilization $U = \frac{\sum_{n=1}^{\infty} \mu_n p_n}{\max_n \mu_n}$

4) Average Response time: $R = \frac{Q}{X}$ Little's law



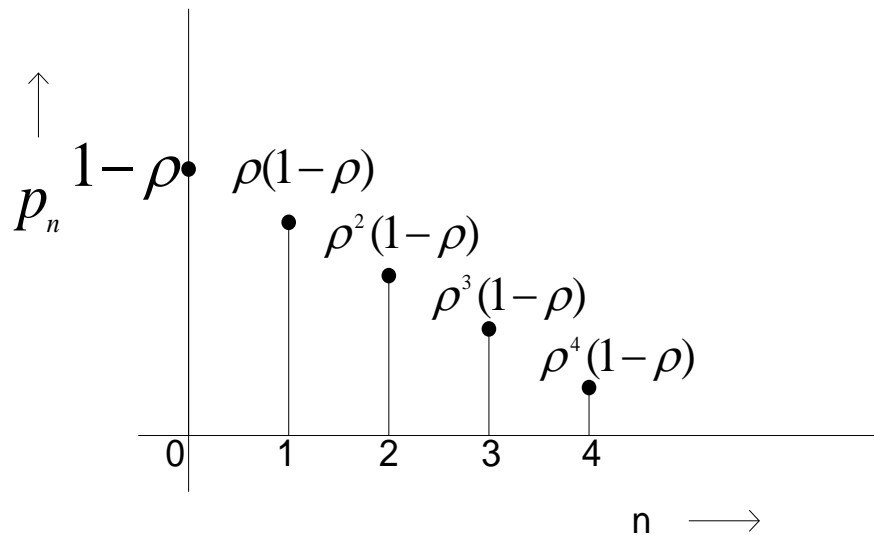
Application to M|M|1 Queue - 1

1. M/M/1 Queue

$$\lambda_n = \lambda, \mu_n = \mu \Rightarrow p_n = \left(\frac{\lambda}{\mu}\right)^n p_0; \frac{\lambda}{\mu} = \rho = \text{traffic intensity}$$

$$p_0 = \left(1 + \sum_{n=1}^{\infty} \rho^n\right)^{-1} = 1 - \rho = 1 - \frac{\lambda}{\mu}; \text{Need } \lambda < \mu \text{ or } \rho < 1$$

$$\therefore p_n = (1 - \rho)\rho^n \quad n = 0, 1, 2, \dots \quad \text{Modified Geometric pmf... note } n \text{ versus } (n - 1)$$





Application to M|M|1 Queue - 2

$$G(z) = \sum_{n=0}^{\infty} p_n z^n = \frac{(1-\rho)}{1-\rho z}$$

$$E(n) = Q = \left. \frac{dG(z)}{dz} \right|_{z=1} = \left. \frac{(1-\rho)\rho}{(1-\rho z)^2} \right|_{z=1} = \frac{\rho}{1-\rho}$$

$$E(n^2) - Q = \left. \frac{d^2G(z)}{dz^2} \right|_{z=1} = \frac{2(1-\rho)\rho^2}{(1-\rho)^3} = \frac{2\rho^2}{(1-\rho)^2}$$

$$\Rightarrow \sigma_n^2 = \frac{2\rho^2}{(1-\rho)^2} + \frac{\rho}{1-\rho} - \left(\frac{\rho}{1-\rho} \right)^2 = \frac{\rho}{(1-\rho)^2} \Rightarrow \sigma_n = \frac{\sqrt{\rho}}{1-\rho}$$

$$C_n = \frac{\sigma_n}{Q} = \frac{1}{\sqrt{\rho}} \geq 1$$

- Throughput $X = \sum_{n=1}^{\infty} \mu_n p_n = \mu \left(\sum_{n=1}^{\infty} \rho^n \right) (1-\rho) = \rho\mu = \lambda$

- Average response time $R = \frac{Q}{X} = \frac{1}{\mu(1-\rho)}$

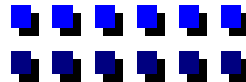
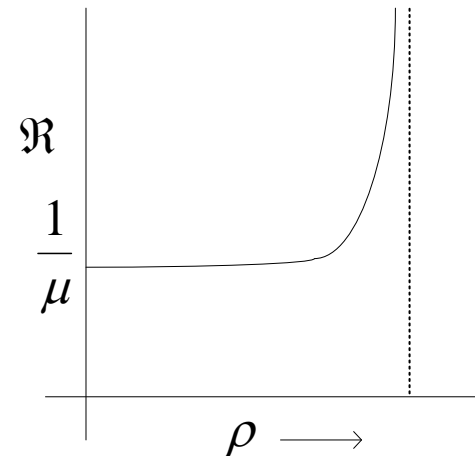
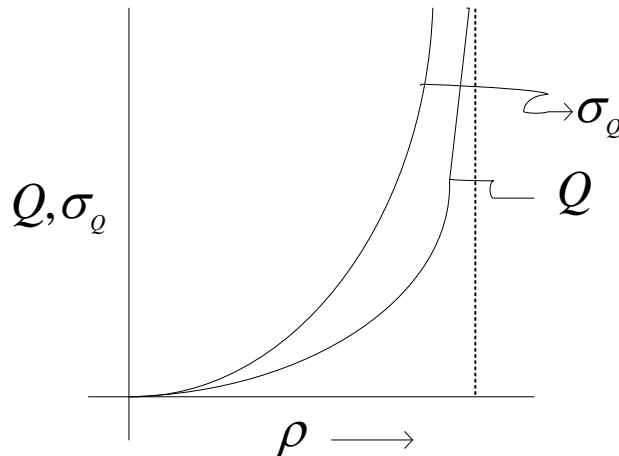
Higher moments of response time requires analysis... later



Application to M|M|1 Queue - 3

- Utilization $U = \frac{\lambda}{\mu} = \rho$
- Average waiting time $W = R - \frac{1}{\mu} = \frac{1}{\mu} \frac{\rho}{1-\rho}$
- Average waiting queue length

$$Q_w = XW = \frac{\rho^2}{1-\rho} \text{ (check: } Q_w = Q - \rho = \frac{\rho}{1-\rho} - \rho = \frac{\rho^2}{1-\rho} \text{)}$$





Application to M|M|∞ Queue

2. M/M/∞ Queue $\lambda_n = \lambda, \mu_n = n\mu$

$$p_n = \frac{\lambda}{n\mu} p_{n-1} \Rightarrow p_n = \frac{\rho^n}{n!} p_0; \quad \rho = \frac{\lambda}{\mu} \Rightarrow p_0 = \left(1 + \sum_{n=1}^{\infty} \frac{\rho^n}{n!}\right)^{-1} = e^{-\rho}$$

The result: $p_n = \frac{\rho^n}{n!} e^{-\rho}; n = 0, 1, 2, \dots$ **POISSON DISTRIBUTION WITH MEAN ρ**

$$\left. \begin{array}{l} Q = \rho \\ \sigma_n = \sqrt{\rho} \end{array} \right\} \Rightarrow C_n = \frac{1}{\sqrt{\rho}}$$

Throughput = $X = \sum_{n=1}^{\infty} \mu_n p_n = \sum_{n=1}^{\infty} \frac{n\mu \cdot \rho^n}{n!} e^{-\rho} = \mu\rho \left(\sum_{n=1}^{\infty} \frac{\rho^{n-1}}{(n-1)!}\right) e^{-\rho} = \lambda$ as it should

$$R = \frac{1}{\mu}$$

$$W = R - \frac{1}{\mu} = 0$$

$$Q_w = 0$$

$$U = \lim_{n \rightarrow \infty} \frac{\lambda}{n\mu} = 0$$

Results valid for M/G/∞ queue



Application to M|M|m Queue -1

1. M/M/m Queue

$$p_n = \frac{\lambda}{n\mu} p_{n-1} \quad \text{for } 1 \leq n \leq m \Rightarrow p_n = \prod_{i=0}^{n-1} \left(\frac{\lambda}{(i+1)\mu} \right) p_0 = \left(\frac{\lambda}{\mu} \right)^n \frac{1}{n!} p_0 \quad 1 \leq n \leq m$$

$$\text{For } n > m: p_n = \frac{\lambda}{m\mu} \cdot p_{n-1} \Rightarrow p_n = \left(\frac{\lambda}{m\mu} \right)^{n-m} p_m; n > m$$

$$\text{Result: } p_n = \begin{cases} \left(\frac{\lambda}{\mu} \right)^n \frac{1}{n!} p_0 & 0 \leq n \leq m \\ \left(\frac{\lambda}{\mu} \right)^m \left(\frac{\lambda}{m\mu} \right)^{n-m} p_0 & n > m \end{cases} \quad \text{From } \sum_{n=0}^{\infty} p_n = 1, \left[\sum_{n=0}^{m-1} \frac{\left(\frac{\lambda}{\mu} \right)^n}{n!} + \sum_{n=m}^{\infty} \frac{\left(\frac{\lambda}{\mu} \right)^m}{m!} \left(\frac{\lambda}{m\mu} \right)^{n-m} \right] p_0 = 1$$

$$\Rightarrow p_0 = \left[\sum_{n=0}^{m-1} \frac{\left(\frac{\lambda}{\mu} \right)^n}{n!} + \frac{\left(\frac{\lambda}{\mu} \right)^m}{m!} \frac{1}{\left(1 - \frac{\lambda}{m\mu} \right)} \right]^{-1}$$



Application to M|M|m Queue - 2

□ Probability of Queuing, P_Q

$$P_Q = \sum_{n=m}^{\infty} p_n = \frac{\left(\frac{\lambda}{\mu}\right)^m}{m!} p_0 \sum_{n=m}^{\infty} \left(\frac{\lambda}{m\mu}\right)^{n-m} = \frac{P_m}{1 - \frac{\lambda}{m\mu}} = \frac{P_m}{1 - \rho}; \rho = \frac{\lambda}{m\mu}$$

(Erlang's delay formula or Erlang's C - formula or Erlang's second formula)

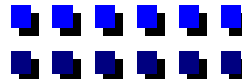
□ Average Number of Customers Waiting (not in service)

$$Q_W = \sum_{n=0}^{\infty} n p_{n+m} = \sum_{n=0}^{\infty} n \frac{\left(\frac{\lambda}{\mu}\right)^m}{m!} \rho^n p_0 = p_m \sum_{n=0}^{\infty} n \rho^n = p_m \rho \cdot \frac{d}{d\rho} \left(\frac{1}{1-\rho} \right) = \frac{p_m \rho}{(1-\rho)^2} = \frac{\rho P_Q}{(1-\rho)}$$

$$\Rightarrow \frac{Q_W}{P_Q} = \frac{\rho}{1-\rho}; \rho = \frac{\lambda}{m\mu}$$

□ Throughput

$$X = \sum_{n=1}^{\infty} p_n \mu_n = m\mu \sum_{n=m}^{\infty} p_n + \sum_{n=1}^{m-1} \frac{\left(\frac{\lambda}{\mu}\right)^n}{n!} \cdot n\mu \cdot p_0 = \lambda \sum_{n=m}^{\infty} p_{n-1} + \lambda \sum_{n=0}^{m-2} p_n = \lambda \text{ as it should}$$





Application to M|M|m Queue - 3

- ❑ Average Waiting Time $W = \frac{Q_w}{\lambda} = \frac{1}{\mu} \frac{P_o}{1-\rho}$
- ❑ Average Response Time $R = W + \frac{1}{\mu} = \frac{1}{\mu} [1 + \frac{P_o}{1-\rho}]$
- ❑ Average Queue Length $Q = R \cdot \lambda = \frac{\lambda}{\mu} [1 + \frac{P_o}{1-\rho}]$
- ❑ Average number of busy servers $Q - Q_w = \frac{\lambda}{\mu}$

Note $\frac{\lambda}{\mu}$ can be > 1 but $\frac{\lambda}{m\mu}$ must be < 1

- ❑ An alternative expression for system length ... a trick that will be useful for state-dependent nodes of networks

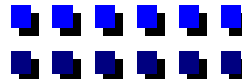
$$p_n = \frac{\lambda}{\mu(n)} p_{n-1}$$

Multi-server

$$\mu(n) = \begin{cases} n\mu & n \leq m \\ m\mu & n > m \end{cases}$$

State-dependent node

$$\{\mu(1) \dots \mu(m)\}$$





Application to M|M|m Queue - 4

Queue Length

$$Q = \sum_{n=1}^{\infty} np_n = \lambda \sum_{n=1}^{\infty} \frac{np_{n-1}}{\mu(n)} = \lambda \left\{ \sum_{n=1}^{m-1} \left[\frac{np_{n-1}}{\mu(n)} - \frac{np_{n-1}}{\mu(m)} \right] + \sum_{n=1}^{\infty} \frac{np_{n-1}}{\mu(m)} \right\}$$

$$Q = \frac{\lambda}{\mu(m)} \left[1 + Q + \sum_{n=1}^{m-1} n \left[\frac{\mu(m)}{\mu(n)} - 1 \right] p_{n-1} \right]$$

or $Q = \left(\frac{\rho}{1-\rho} \right) [1 + \gamma]$ for multi-server nodes
 $\gamma = 0$ for single server nodes

$\gamma = \sum_{n=1}^{m-1} (m-n)p_{n-1}$ useful for numerical computations

$\rho = \frac{\lambda}{m\mu}$; Note: need p_0, p_1, \dots, p_{m-2} only

Utilization $U = \frac{\lambda}{m\mu}$



Higher Moments for M|M|1 Queue

Higher Moments of Queue Length

$$\begin{aligned} E(n^k) &= \sum_{n=1}^{\infty} n^k p_n = \rho \sum_{n=1}^{\infty} n^k p_{n-1} = \rho \sum_{n=1}^{\infty} (n-1+1)^k p_{n-1} = \rho \sum_{n=1}^{\infty} \sum_{r=0}^k \binom{k}{r} (n-1)^r p_{n-1} \\ &= \rho \sum_{r=0}^k \binom{k}{r} E[n^r] \end{aligned}$$

$$\Rightarrow E(n^k) = \frac{\rho}{1-\rho} \left\{ \sum_{r=0}^{k-1} \binom{k}{r} E[n^r] \right\}$$

$$k=1 \Rightarrow Q = \frac{\rho}{1-\rho}$$

$$k=2 \Rightarrow E(n^2) = \frac{\rho}{1-\rho} (1+2Q) = \frac{\rho(1+\rho)}{(1-\rho)^2} \Rightarrow \sigma_n^2 = \frac{\rho}{(1-\rho)^2} = \frac{Q}{(1-\rho)} \Rightarrow C_n = \frac{\sigma_n}{Q} = \frac{1}{\sqrt{\rho}}$$

$$k=3 \Rightarrow E(n^3) = \frac{\rho}{1-\rho} [1+3Q+3E(n^2)] = \frac{\rho(1+4\rho+\rho^2)}{(1-\rho)^3}$$

$$\Rightarrow \text{skewness} = E[(n-Q)^3] = \frac{\rho(1+\rho)}{(1-\rho)^3} = \frac{E(n^2)}{(1-\rho)}$$



Summary of Lecture 3

- Properties of Exponential and Geometric Distributions
- Properties of the Poisson Process
- Steady-state solutions to DTMC and CTMC
- State Transition Rate Diagrams
- Applications to $M/M/1$, $M/M/1/N$, $M/M/m$, $M/M/\infty$, $M/M/m/m$ Queues
- Performance measures for $M/M/1$, $M/M/\infty$, and $M/M/m$ Queues